

Embedding Lemma

We now introduce a generalization of the triangle counting lemma. There are many versions and strengthenings of this lemma often called the counting lemma or embedding lemma.

Lemma 1 (Improved embedding lemma). *Let F be a k -chromatic graph with maximum degree $\Delta(F)$. Fix $0 < \delta < \frac{1}{k}$ and let G be a graph and let V_1, \dots, V_k be disjoint sets of vertices of G . If each V_i has $|V_i| \geq 2\delta^{-\Delta}|F|$ and each pair of partition classes is $\frac{1}{2\Delta}\delta^\Delta$ -regular with density $\geq 2\delta$, then the classes V_1, V_2, \dots, V_k contain F .*

Proof. □

Lemma 2 (“Weak” embedding lemma). *Let F be a k -chromatic graph on f vertices ($2 \leq k \leq f$). Fix $0 < \delta < \frac{1}{k}$ and let G be a graph and let V_1, \dots, V_k be disjoint sets of vertices of G . If each V_i has $|V_i| \geq \delta^{-f}$ and each pair of partition classes is δ^f -regular with density $\geq 2\delta$, then the classes V_1, V_2, \dots, V_k contain F .*

Proof. We will prove the lemma with the two following strengthenings. First, we only require that the pairs of V_1, \dots, V_k have density at least $2\delta - \delta^f$. Second, we will show that if we (arbitrarily) label the color classes of F , then we can embed F into G such that the first color class of F is in V_1 and all other classes are in $V_2 \cup V_3 \cup \dots \cup V_k$.

We proceed by induction on f . For $f = 2$ the statement is trivial. Let $f > 2$ and assume the theorem holds for smaller values. Given F , let us remove a vertex y in the first color class to get $F' = F - y$.

For $2 \leq i \leq k$, let R_i be the set of vertices in V_1 that are adjacent to less than $\delta|V_i|$ vertices of V_i .

1: Show that $V_1 - \cup R_i$ is not empty. Hint: Use regularity of V_1, V_i to show that R_i is not too big.

Solution: If $|R_i| \geq \delta^f|V_1|$, then we have $|d(V_1, V_i) - d(R_i, V_i)| < \delta^f$ which implies that $d(V_1, V_i) < d(R_i, V_i) + \delta^f < \delta + \delta^f < 2\delta - \delta^f$; a contradiction. Thus, we have $|R_i| < \delta^f|V_1|$, so $|\cup R_i| < (k-1)\delta^f|V_1| < |V_1|$.

Thus, there is a vertex $x \in V_1 - \cup R_i$. Put $V'_1 = V_1 - x$ and $V'_i = V_i \cap N(x)$ (for $i \neq 1$).

Now our goal is to show that V'_1, \dots, V'_k satisfy the inductive hypothesis for $f-1$.

2: First verify that V'_i is large enough for all i .

Solution: Observe that

$$|V'_1| = |V_1| - 1 \geq \delta|V_1| \geq \delta^{-(f-1)}$$

and for $i \neq 1$,

$$|V'_i| \geq \delta|V_i| \geq \delta^{-(f-1)}.$$

3: Check that V'_i, V'_j is a regular pair for all i, j .

Solution: Now suppose $i \neq j$ and $A \subset V'_i$ and $B \subset V'_j$ are such that $|A| \geq \delta^{f-1}|V'_i| \geq \delta^f|V_i|$ and $|B| \geq \delta^{f-1}|V'_j| \geq \delta^f|V_j|$, then, with an application of the triangle inequality, we get

$$|d(V'_i, V'_j) - d(A, B)| \leq |d(V'_i, V'_j) - d(V_i, V_j)| + |d(V_i, V_j) - d(A, B)| < 2\delta^f < \delta^{f-1}.$$

4: Check that $d(V'_i, V'_j)$ is large enough for all i, j .

Solution: Finally,

$$|d(V_i, V_j) - d(V'_i, V'_j)| < \delta^f$$

implies that

$$d(V'_i, V'_j) > d(V_i, V_j) - \delta^f \geq 2\delta - 2\delta^f > 2\delta - \delta^{f-1}.$$

Thus V'_1, V'_2, \dots, V'_k satisfy the conditions of the theorem for $f - 1$.

5: Use induction to finish the proof.

Solution: By induction we can embed F' into $G - x$ such that the first color class of F' is in V'_1 and the other color classes of F' are in V'_2, \dots, V'_k . Because $x \in G$ is adjacent to every vertex in V'_2, \dots, V'_k we can put $y = x$ to embed F into G as desired. □

We are now ready to give another proof of Erdős-Stone-Siminovits Theorem 3) using the regularity lemma and embedding lemma. Recall

Theorem 3. Let F be a graph with chromatic number $\chi(F)$ and fix $0 < \delta < \frac{1}{\chi(F)}$, then there is an $n_0 = n_0(\delta, |F|)$ such that if G is a graph on $n \geq n_0$ vertices and

$$e(G) > \left(1 - \frac{1}{\chi(F) - 1} + \delta\right) \frac{n^2}{2},$$

then G contains F as a subgraph.

Proof. Put $f = |F|$ and let G be as in the statement of the theorem. Let us apply the regularity lemma to G to with $\epsilon = (\frac{\delta}{8})^f$ and $m > \frac{8}{\delta} > \chi(F)$. That is, there is an equipartition V_1, \dots, V_r of G into r parts such that $\frac{8}{\delta} < r < M$ and all but at most $(\frac{\delta}{8})^f r^2$ of the pairs of clusters are $(\frac{\delta}{8})^f$ -regular.

6: As in triangle removal lemma, count the number of edges 1. inside each cluster, 2. edges between clusters that are not regular, 3. edges between clusters with $d(V_i, V_j) < \frac{\delta}{4}$.

Solution: As in the triangle removal lemma we will remove the following edges.

1. Remove the edges inside of each cluster V_i . There are at most $r \binom{\lceil n/r \rceil}{2} \leq \frac{n^2}{r} < \frac{\delta}{8} n^2$ such edges.
2. Remove the edges between all pairs V_i, V_j that are not δ^f -regular. There are at most $2(\frac{\delta}{8})^f r^2$ such pairs and each has at most $(\frac{n}{r})^2$ edges. So we remove at most $(\frac{\delta}{8})^f n^2 \leq \frac{\delta}{8} n^2$ such edges
3. Remove the edges between all pairs V_i, V_j if the density of the pair $d(V_i, V_j) < \frac{\delta}{4}$. There are less than $\binom{r}{2} \frac{\delta}{4} (\frac{n}{r})^2 < \frac{\delta}{8} n^2$ such edges.

In total we have removed at less than $\delta \frac{n^2}{2}$ edges, so we still have more than

$$\left(1 - \frac{1}{\chi(H) - 1}\right) \frac{n^2}{2}$$

edges.

7: Use Turán's theorem and the embedding lemma to finish the proof.

Solution: So, by Turán's theorem, the remaining graph contains a clique on $\chi(F)$ vertices. The $\chi(F)$ clusters that contain the vertices of this clique satisfy the conditions of the embedding lemma with $\frac{\delta}{8}$, so G contains F . □

Now we go back and show application of the triangle removal lemma. Here we restate it for convenience.

Theorem 4 (Triangle removal lemma, Ruzsa-Szemerédi, 1978). *For $\alpha > 0$, there exists $\beta > 0$ such that if G is an n -vertex graph that requires the removal of αn^2 edges to be triangle-free, then G has at least βn^3 triangles.*

Recall that a k -term **arithmetic progression** (AP) is a sequence a_1, a_2, \dots, a_k such that there is some fixed d such that $a_{i+1} - a_i = d$ for all i .

Theorem 5 (Roth's theorem, 1953). *If $S \subset [n]$ contains no 3-term arithmetic progression, then $|S| = o(n)$.*

Proof. (Ruzsa-Szemerédi) Fix $\epsilon > 0$, and let $S \subset [n]$ of size ϵn . We will show that for all n large enough, that S contains a 3-term AP. We construct a graph G with vertex set partitioned into three classes $A = [n], B = [2n], C = [3n]$. If $s \in S$ and $x \in [n]$, then add to G the triangle between vertices $x \in A, x + s \in B$, and $x + 2s \in C$. Observe that for $b \in B$ and $c \in C$ forming an edge we necessarily have $c - b \in S$.

8: Use Triangle removal lemma to show that there are other triangles than just the ones we added. Hint: How many we added?

Solution: Clearly, the $n|S| = \epsilon n^2$ triangles formed in this way are edge-disjoint. To destroy these edge-disjoint triangles, we must remove ϵn^2 edges of G . Therefore, by the triangle removal lemma there are at least δn^3 triangles. For n large enough (depending on ϵ) we have $\delta n^3 > \epsilon n^2$ which implies that G contains a triangle other than those explicitly defined above.

9: Let x be the vertex of the triangle in A , that is not one of the explicitly added ones. Describe the triangle and use it to find a 3-term AP.

Solution: Because this triangle is different from those defined above we have distinct $b, c \in S$ such that $x + b \in B$ and $x + 2c \in C$ are the other vertices of the triangle. Thus $(x + 2c) - (x + b) = 2c - b$ is necessarily an element a of S . That is, $c - b = a - c$, i.e., b, c, a is a 3-term AP. □

Theorem 6 (Solymosi, 2001). *For any $\alpha > 0$, there exists N such that if A is a set of $\geq \alpha N^2$ many points on the $N \times N$ integer lattice, then A contains three distinct points of the following form $(x, y), (x + d, y), (x, y + d)$, i.e., an isosceles right triangle.*

Proof. Consider the collection of vertical lines, horizontal lines and 45° diagonal (left to right) lines in the $N \times N$ lattice.

Construct a 3-partite graph G with classes X, Y, Z such that X is the set of vertical lines, Y is the set of horizontal lines, and Z is the set of diagonal lines.

Two vertices (lines) in this graph are connected by an edge if the intersection of the two lines is an element of A . Therefore, for each element $a \in A$, the three lines intersecting in a form a triangle in G .

10: Finish the proof by looking at triangles. □